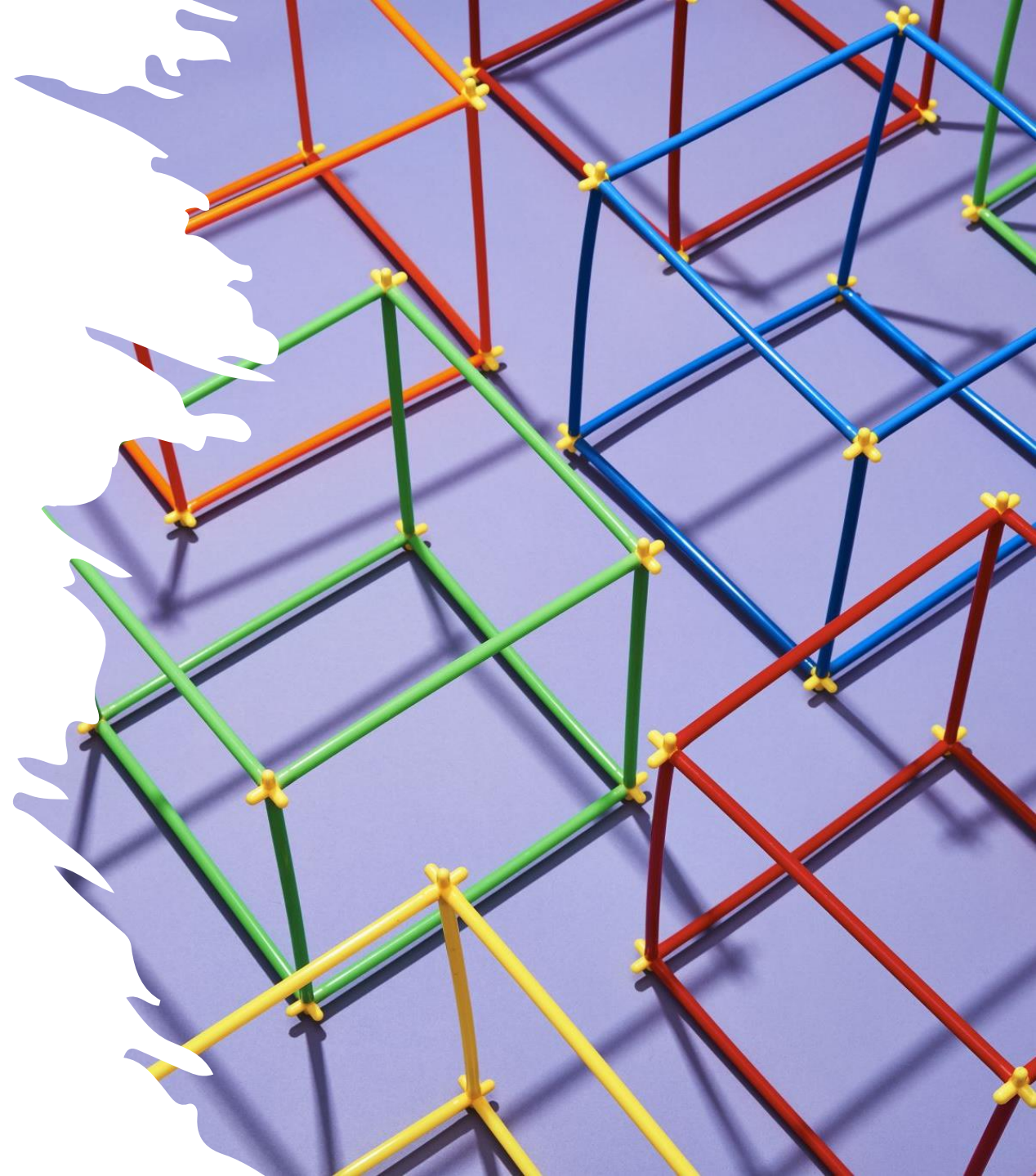


Data Structure

Lecture 4: Linked List

Prepared by

Dr. Mohammed Salah Al-Obiadi

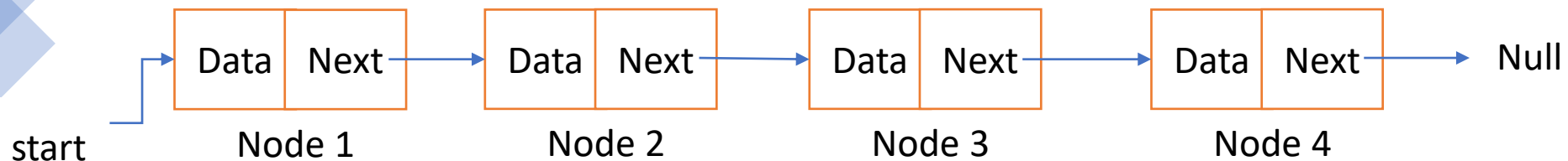


What is a Linked List?

A linked list is a data structure used for storing collections of data.

A linked list has the following properties:

1. Successive elements are connected by pointers.
2. The last element points to NULL.
3. Can grow or shrink in size during execution of a program.
4. Can be made just as long as required (until systems memory exhausts).
5. Does not waste memory space. It allocates memory as list grows.



Linked List vs Arrays?

Array	Linked list
Array elements store in <u>a contiguous memory location</u> .	Linked list elements can be stored <u>anywhere in the memory</u>
Array works with <u>a static memory</u> and cannot be changed at the run time.	The Linked list works with <u>dynamic memory</u> means memory size can be changed at the run time.
Array elements are <u>independent</u> of each other.	Linked list elements are <u>dependent</u> on each other. As each node contains the address of the next node.
Array takes <u>more time</u> while performing any operation like insertion, deletion, etc.	Linked list takes <u>less time</u> while performing any operation like insertion, deletion, etc.
Accessing any element in an array is <u>faster</u> as the element in an array can be directly accessed through the index.	Accessing an element in a linked list is <u>slower</u> as it starts traversing from the first element of the linked list.
In the case of an array, memory is allocated at <u>compile-time</u> .	In the case of a linked list, memory is allocated at <u>run time</u> .
Memory utilization is <u>inefficient</u> in the array. For example, if the size of the array is 6, and array consists of 3 elements then the rest of the space will be unused.	Memory utilization is <u>efficient</u> as the memory can be allocated or deallocated at the run time.
Arrays take $O(1)$ for access to an element.	Linked lists take $O(n)$ for access to an element.

Operation on Linked List

1- Traversal: To traverse all the nodes one after another.

2- Insertion: To add a node at the given position.

3- Deletion: To delete a node.

4- Searching: To search an element(s) by value.

5- Updating: To update a node.

6- Sorting: To arrange nodes in a linked list in a specific order.

7- Merging: To merge two linked lists into one.

Types of Link List

1- Single Link List

2- Double Link List

3- Circular Link List

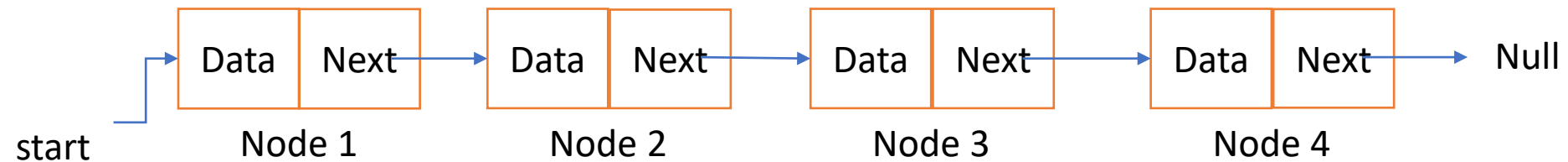
4- Doubly Circular linked list

Single Link List

Generally “linked list” means a single linked list.

This list consists of a number of nodes in which each node has a *next* pointer to the following element.

The link of the last node in the list is NULL, which indicates the end of the list.



STRUCTURE OF THE NODE OF A LINKED LIST

Struct **tagname**

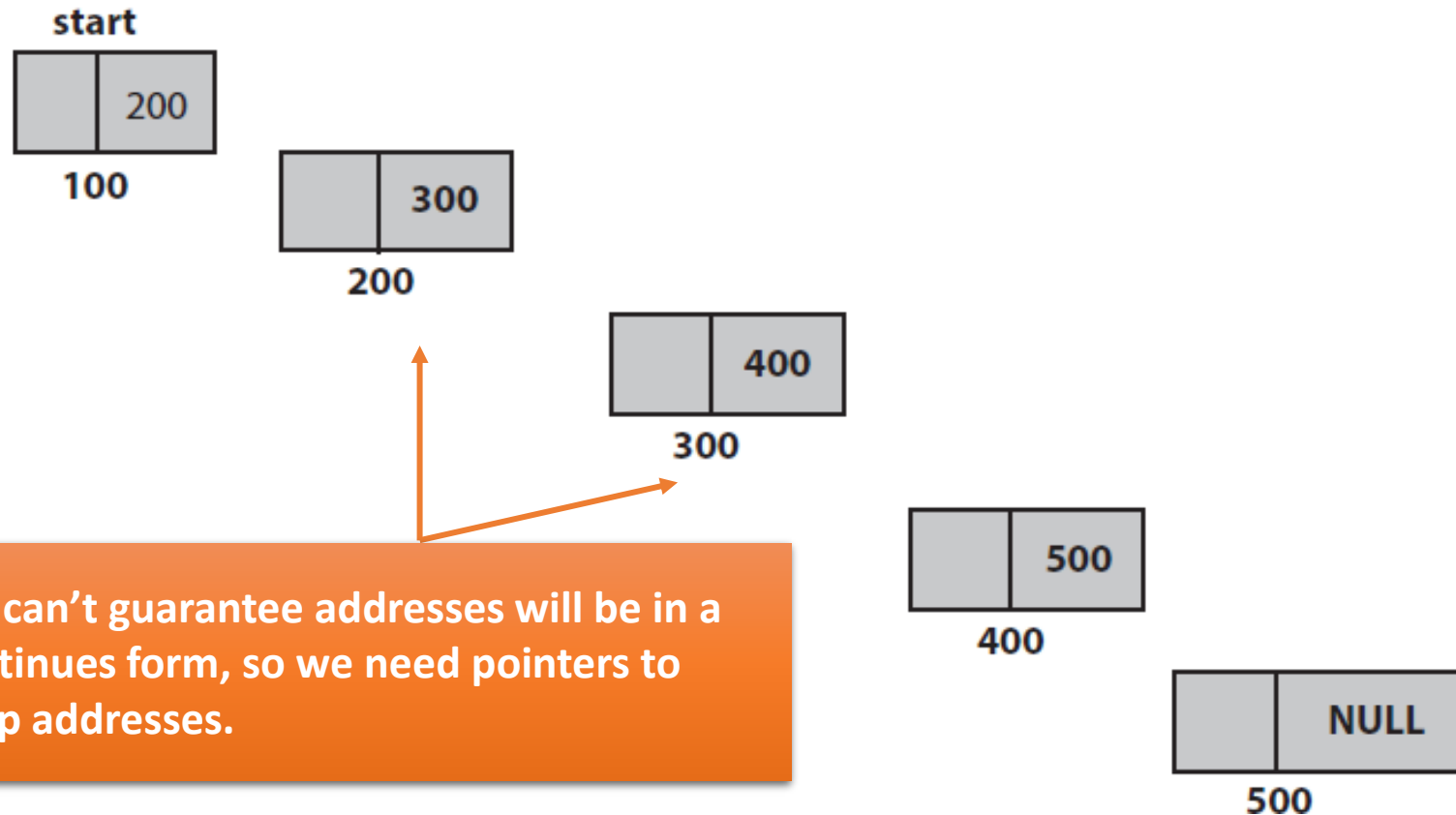
```
{  
    Data type member1;  
    Data type member2;  
    .....  
    .....  
    .....  
    Data type membern;  
    Struct tagname *var;  
};
```

Example:

```
struct link  
{  
    int info;  
    struct link *next;  
};
```

LOGIC FOR CREATION

```
struct link start, *node;
```



Algorithm For Creation Of Single Link List

Struct link start, *node

create(start,node) [start is the structure type of variable][node is the structure type of pointer]

step-1 : node = &start

step-2 : node → next = new link() //allocate memory of size struct link for the node

node = node → next

input : node → info

node → next = null

step-3 : repeat step-2 to create more nodes

step-4 : return



Algorithm For Traversing Of Single Link List

struct link start, *node;

traverse(start,node) [start is the structure type of variable] [node is the structure type of pointer]

step-1 : node = start.next

step-2 : repeat while (node!=null)

 write : node → info

 node = node → next

 end of loop

step-3 : return



Insertion Into Linked List

The insertion process with link list can be discussed in four different ways:

1. Insertion at Beginning.
2. Insertion at End.
3. Insertion when node number is known.
4. Insertion when information is known.

Algorithm For Insertion At Beginning

struct start, *first, *node,* newnode

insbeg(start,first,node, newnode) [start is the structure variable] [node and first is the structure pointer]

step-1 : first = &start //first saves start's address

node = start.next

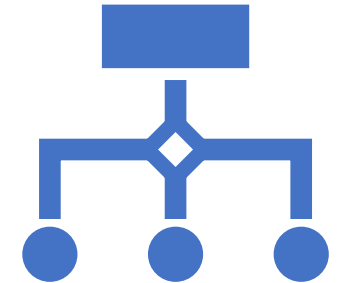
step-2 : newnode = new link()

input : newnode → info

first → next = newnode

newnode → next := node

step-3 : return



Algorithm For Insertion At Last

struct start, *last, *node,* newnode

inslast(start,last,node,newnode)

step-1 : last = &start //last's pointer saves start's address

node = start.next

step-2 : repeat while(node != null)

node = node → next

last = last → next

step-3 : newnode →next=new link() //allocate a memory to newnode

input : newnode → info

last → next = newnode

newnode → next = null

step-4 : return

